

Mon, Apr 29, 2024

* HW 6 due on Friday

* Course Evals open

Topic 13 - Tabu Search

Pseudocode:

generation = 0

taboo = dict()

taboo_time = 20 # depends on the problem

x = random elt. of search space

while True:

generation += 1

neighbors = nbhd(x)

each neighbor is a pair (s,m) where s is the new sol (the neighbor) and m is the "move" that turned $x \rightarrow s$

new_x, move = the pair (s,m) in neighbors with the highest score subject to the constraint that either

m is not a key in taboo
OR $\text{taboo}[m] \leq \text{generation}$

$x = \text{new}_x$

$\text{taboo}[\text{move}] = \text{generation} + \text{taboo_time}$

Advanced Topics

* "Aspiration Criteria" You can decide ahead of time to ignore the taboo list in some cases, for example if a taboo solution would be a new best record.

* Sometimes the neighborhoods are too big. You can try to generate a fixed # of random tweaks (100? 1000?) and pretend that's the whole neighborhood.

* Add extra exploration or exploitation depending on the performance of the algorithm.

- Keep track of how often components of solutions are being used in good or bad or any candidate.

Ex: How often a particular item is used in the 100 best or 100 worst knapsack solutions.

If you want more exploitation:
Force the items in the best sols to stay.

If you want more exploration:
Ban the good items.

Topic 14 - Neighborhoods in Continuous Space

So far we've used a simple tweak in our continuous problems.

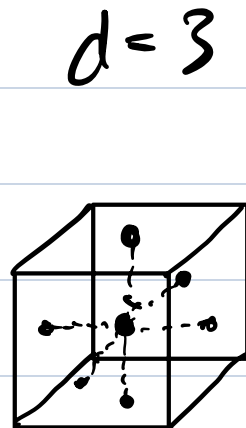
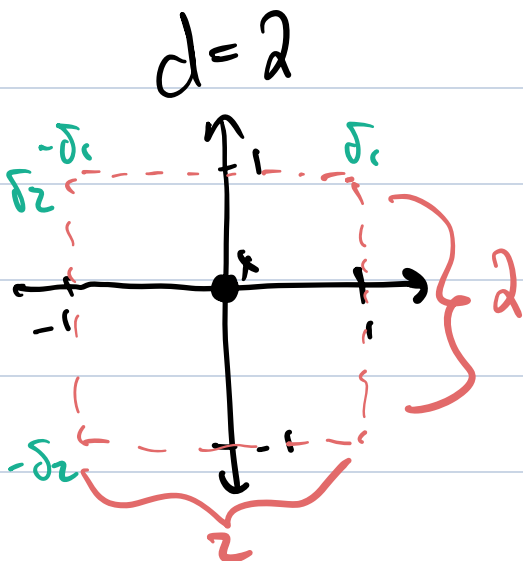
$$x = (x_1, x_2, \dots, x_d) \quad d \text{ dimensions}$$

$s = \text{tweak}(x) = x + (r_1 \delta_1, r_2 \delta_2, \dots, r_d \delta_d)$
 where each r_i is a uniform random # in $[-1, 1]$ and each δ_i is decided ahead of time.

Works okay when d is small (2 or 3)

For now, assume $\delta_i = 1$ and
 $x = (0, 0, 0, \dots, 0)$ so
 $\text{tweak}(x) = (r_1, r_2, \dots, r_d)$.

The new point $\text{tweak}(x)$ is somewhere in the d -dimensional cube with side length 2 centered at the origin.



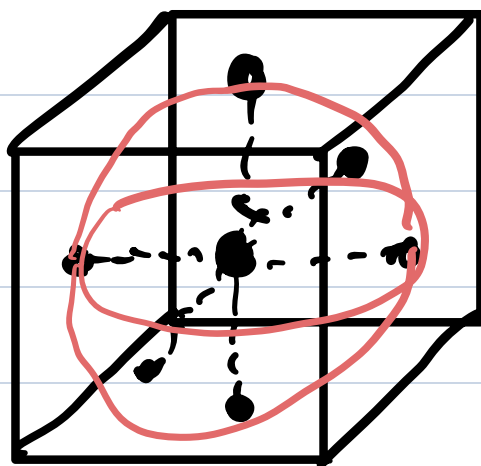
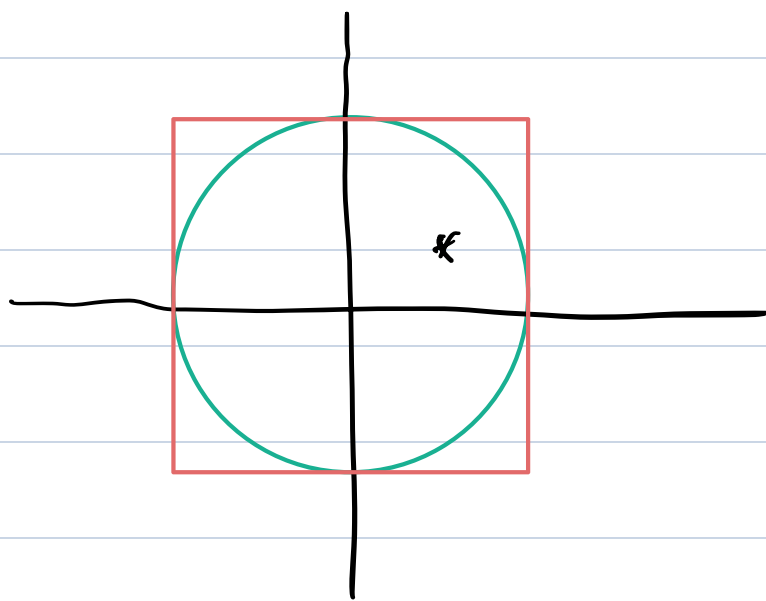
What is the farthest that s could be away from x ?

$$d=2 \quad \leadsto \quad \sqrt{2} \approx 1.41$$

$$d=3 \quad \quad \quad \sqrt{3} \approx 1.7$$

In d -dimensional space $\leadsto \sqrt{d}$

Picking points in a cube can lead to very far away tweaks. Instead pick from a sphere.



* How do you pick points uniformly from the unit circle?

Idea 1) Rejection Sampling

Pick points in the square

Check if each point is in the circle

If so, keep it

If not, throw away and try again

Works, but very slow in higher dimensions

Idea 2) Pick $x \in [-1, 1]$

Pick $y \in [-\sqrt{1-x^2}, +\sqrt{1-x^2}]$