

## LECTURE 22 – THE STANDARD FUNCTION SCALE, PART 1

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There are a number of topics that we're skipping over, but I want you to know that they're in the *Analytic Combinatorics* book in case you find that you need them some day.

- (1) Periodic and non-periodic fluctuations, conditions for proving periodicity (the Dafodil Lemma).
- (2) Localization of zeros and poles, winding number, the Argument Principle, Rouché's Theorem.
- (3) Derivation of exponential growth rate from functional equation (e.g.,  $f(z) = ze^{f(z)}$ ,  $f(z) = z + f(z^2 + z^3)$ ,  $f(z) = (1 - zf(z^2))^{-1}$ )

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In this section we begin to move past the simple realm of polar singularities into more general territory. Still the basic principle holds that *the asymptotic expansion of a function near its singularities* and *the asymptotic behavior of the function's coefficients* are deeply intertwined.

Our goal, now, is to analyze the asymptotic behavior of sequences coming from a broader class of generating functions with “algebraic-logarithmic” singularities. To do this, we need to central types of theorems.

The first type are *standard function scale* results that let us translate the dominant terms in the asymptotic expansion of a function to the dominant terms in the asymptotic behavior of the sequence, while the second type are *transfer theorems* that let us translate the corresponding error terms.

For example<sup>1</sup>, consider a function  $f(z)$  such that

$$f(z) \underset{z \rightarrow 1}{\sim} \frac{1}{(1-z)^{3/2}} + O\left((1-z)^{-1/2}\right).$$

We shall see that if  $\{f_n\}$  is the power series expansion of  $f(z)$  at the origin, then the asymptotic expansion above begets the following asymptotic behavior if  $\{f_n\}$ :

$$f_n \sim n^{1/2} + O(n^{-1/2}).$$

The correspondence between the first terms comes from the standard function scale, while the correspondence between the second terms comes from a transfer theorem.

### BIG-O NOTATION

Big-O notation (and its variants—little-o, big-Omega, etc) give a concise way to describe the growth of a function of a sequence.

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<sup>1</sup>Postponing just one more time the formal definition of  $O(\cdot)$ .

Loosely, think of Big-O as  $\leq$ , and little-o as  $<$ . Formally, we say

$$f(z) = O(g(z))$$

as  $z \rightarrow z_0$  if there is a positive constant  $M$  and a neighborhood of  $z_0$  such that

$$|f(z)| \leq M|g(z)|$$

for all  $z$  in the neighborhood. When  $z_0 = \infty$ , change “neighborhood of  $z_0$ ” to “sufficiently large”. In the same way, we say

$$a(n) = O(b(n))$$

as  $n \rightarrow \infty$  if there exists a positive constant  $M$  such that for  $n$  sufficiently large

$$|a(n)| \leq M|b(n)|.$$

When the second function ( $g(z)$  or  $b(n)$ ) is not zero in the limit, these conditions are equivalent to

$$\limsup_{z \rightarrow z_0} \left| \frac{f(z)}{g(z)} \right| < \infty \quad \text{and} \quad \limsup_{n \rightarrow \infty} \left| \frac{f(n)}{g(n)} \right| < \infty.$$

Really though,  $O(g(z))$  and  $O(b(n))$  are *equivalence classes*, and we should be saying  $f(z) \in O(g(z))$ —but *c’est la vie*.

While  $f(z) = O(g(z))$  roughly means “ $f(z)$  is on the order of  $g(z)$  or smaller”, little-o notation is stronger. When one write  $f(z) = o(g(z))$ , they mean “ $f(z)$  is of a strictly smaller order than  $g(z)$ ”. Formally,  $f(z) = o(g(z))$  as  $z \rightarrow z_0$  if

$$\lim_{z \rightarrow z_0} \frac{f(z)}{g(z)} = 0.$$

Similar adjustments account for  $z \rightarrow \infty$  and for the case of sequences.

There are other asymptotic measures as well.

- We say  $f(z) = \Omega(g(z))$  if  $f(z)$  is asymptotically at least as large as  $g(z)$  (up to constant multipliers).
- We say  $f(z) = \omega(g(z))$  if  $f(z)$  is asymptotically strictly larger than  $g(z)$ .
- We say  $f(z) = \Theta(g(z))$  if  $f(z)$  is asymptotically exactly as large as  $g(z)$  (up to constant multipliers).

### Examples:

- As  $x \rightarrow 0$ ,  $e^x = 1 + x + O(x^2)$
- As  $x \rightarrow 0$ ,  $e^x = 1 + x + o(x^{3/2})$
- As  $n \rightarrow \infty$ ,  $n^3 + 2n^2 + 1 = O(n^{100})$
- As  $n \rightarrow \infty$ ,  $n^3 + 2n^2 + 1 = O(n^3)$
- As  $n \rightarrow \infty$ ,  $n^3 + 2n^2 + 1 = o(n^{100})$
- As  $n \rightarrow \infty$ ,  $n^3 + 2n^2 + 1 = \Theta(n^3)$
- As  $n \rightarrow \infty$ ,  $n^2 = O(10^{10^{10}} n^2)$  and  $n^2 = \omega(10^{10^{10}} n^{1.999})$

In particular, as a function approaches a singularity, its growth is dominated by the additive component that goes to infinity fastest.

### THE STANDARD FUNCTION SCALE

Our first step is to find exactly the asymptotic expansions of functions whose expansion near a singularity  $\zeta$  has the form

$$f(z) \underset{z \rightarrow \zeta}{\sim} \left(1 - \frac{z}{\zeta}\right)^{-\alpha} \left(\log\left(\frac{1}{1 - \frac{z}{\zeta}}\right)\right)^\beta.$$

We shall see that this uniformly gives rise to sequences with asymptotic behavior

$$f_n \sim C \cdot \zeta^{-n} n^{\alpha-1} \log(n)^B,$$

where  $B = \beta$  or  $B = \beta - 1$ .

As before, we consider only functions whose dominant singularity is at  $z = 1$ , which causes no loss of generality.

**Example:** Before we dive in, we present one more example of the power of the forthcoming methods. The class of labeled 2-regular graphs is constructed by

$$\mathcal{R} = \text{SET}(\text{UCYC}_{\geq 3}(\mathcal{Z}))$$

and so it has the exponential generating function

$$R(z) = \exp\left(\frac{1}{2} \left(\log(1-z)^{-1} - z - \frac{z^2}{2}\right)\right) = \frac{e^{-z/2 - z^2/4}}{\sqrt{1-z}}.$$

First, we note that the expansion of  $e^{-z/2 - z^2/4}$  around  $z = 1$  is

$$e^{-z/2 - z^2/4} = e^{-3/4} - e^{-3/4}(z-1) + \frac{e^{-3/4}}{4}(z-1)^2 + \frac{e^{-3/4}}{12}(z-1)^3 + \dots$$

Therefore, near  $z = 1$ ,

$$e^{-z/2 - z^2/4} = e^{-3/4} + O(1-z).$$

Thus,

$$\frac{e^{-z/2 - z^2/4}}{\sqrt{z-1}} = \frac{e^{-3/4}}{\sqrt{1-z}} + O(\sqrt{1-z}).$$

Our forthcoming work will tell us immediately that

$$R_n \sim \frac{e^{-3/4}}{\sqrt{\pi n}} + O(n^{-3/2}),$$

and more complete asymptotic expansions can be found starting with, for example,

$$e^{-z/2 - z^2/4} = e^{-3/4} - e^{-3/4}(z-1) + \frac{e^{-3/4}}{4}(z-1)^2 + O((1-z)^3).$$

**Polynomial Factors.** We have already made a similar calculation for rational generating functions. If  $f(z) = (1 - z)^{-\alpha}$  for a positive integer  $\alpha$ , then by Newton's Generalized Binomial Theorem,

$$\begin{aligned} [z^n]f(z) &= (-1)^n \binom{-\alpha}{n} \\ &= \binom{n + \alpha - 1}{n} \\ &= \frac{(n + \alpha - 1)(n + \alpha - 2) \cdots (n + 1)}{(\alpha - 1)!} \\ &\sim \frac{n^{\alpha-1}}{(\alpha - 1)!}. \end{aligned}$$

What we seek is a method that works for all  $\alpha$  and gives an asymptotic expansion to arbitrary accuracy. For this we turn to complex analysis.

**Theorem 22.1.** *Let  $\alpha$  be an arbitrary complex number. The coefficient of  $z^n$  in  $f(z) = (1 - z)^{-\alpha}$  admits, for large  $n$ , a complex asymptotic expansion in descending powers of  $n$ ,*

$$[z^n]f(z) \sim \frac{n^{\alpha-1}}{\Gamma(\alpha)} \left( 1 + \sum_{k=1}^{\infty} \frac{e_k(\alpha)}{n^k} \right),$$

where  $e_k(\alpha)$  is a computable polynomial in  $\alpha$  of degree  $2k$ , divisible by  $\alpha(\alpha - 1) \cdots (\alpha - k)$ , and where  $1/\Gamma(\alpha)$  is understood to vanish when  $\alpha \in \mathbb{Z}_{\leq 0}$ . In particular,

$$e_k(\alpha) = \sum_{\ell=k}^{2k} (-1)^\ell \lambda_{k,\ell} (\alpha - 1) \cdots (\alpha - \ell),$$

where  $\lambda_{\ell,k} = [v^k t^\ell] e^t (1 + vt)^{-1-1/v}$ . For example,

$$[z^n]f(z) = \frac{n^{\alpha-1}}{\Gamma(\alpha)} \left( 1 + \frac{\alpha(\alpha - 1)}{2n} + \frac{\alpha(\alpha - 1)(\alpha - 2)(3\alpha - 1)}{24n^2} + O\left(\frac{1}{n^3}\right) \right).$$

*Proof.* The main strategy is to use Cauchy's coefficient formula and integrate around a particular contour (a *Hankel contour*, to be precise), that approaches the singularity at  $z = 1$  closely before traveling away from it. Additionally, a change of variables  $z = 1 + t/n$  and some precise estimates are needed. See the textbook for the gory details.  $\square$