

Monday, Feb. 19, 2024  
Scientific Computing

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## Announcements:

- Office hours canceled today  
Tuesday, 11am, Teams
- HW 2 due on Friday

Topic 5 - Search Spaces + Brute Force  
(continued)

Example: Gameshop problem.

What does a possible solution look like?

You have 60 transaction slots and you need to assign a person to each one. If you start with  $n$  people, how many ways can this be done?

60 slots: \_ \_ \_ \_ \_ . . . . \_

(Not just valid assignments — all assignments)

Slot 1:  $n$  people

Slot 2:  $n-1$  people

Slot 3:  $n-2$

⋮

Slot 60:  $n-59$  people

Total # of configs:

$$n \cdot (n-1) \cdot (n-2) \cdot (n-3) \cdots (n-59) = \frac{n!}{(n-60)!}$$

Search space: all ordered lists of 60 people

Multiply out:  $n^{60} + [\text{stuff with powers less than } 60]$

Size:  $O(n^{60})$

Good news: polynomial, not exponential

Bad news: still pretty bad

10 people: A B C D E F G H I J

3 PS  $\underline{10} \cdot \underline{9} \cdot \underline{8} = 720$

(A, B, C)

(A, C, B)

(B, A, C)

(C, B, A)

(C, A, B)

(B, C, A)

# NFL Scheduling =

search space per week = all ways of putting 32 teams in pairs

For 17 weeks: do this 17 times

$$\approx 6.5 \times 10^{294}$$

(ignoring bye weeks)

$10^{20} \approx$  the # atoms in the

UNIVERSE

Summary: (brute force)

python package called "itertools"

Pros:

very easy to code

fewer bugs

guaranteed optimal

find all optimal solutions

good for testing other methods against

(1) if you're coding a different guaranteed optimal method, check that it works correctly (for small data)

(2) if you're coding a non-guaranteed

optimal algorithm, testing how close does it get to optimal for small data

## Topic 6 - Divide and Conquer (D+C)

D+C is an algorithmic paradigm (a problem solving approach) that roughly goes like:

- 1) Split the input data in half
- 2) Solve the problem on each half separately (recursion!)
- 3) Combine your two answers into one big answer.

Classic Example: Sorting a list

\* Let's phrase this as a constraint satisfaction problem

\* Input: list of  $n$  numbers

\* Search space: All orderings of  $n$  things. These are called "permutations" and the # of them is:

$$n \cdot (n-1) \cdot (n-2) \cdots 3 \cdot 2 \cdot 1$$

first item      2<sup>nd</sup> item      3<sup>rd</sup> item      "

n!

\* Goal: Find the rearrangement that puts things in the right order.

(smallest to largest)

An "obvious" optimal algorithm: (greedy-ish)

- $n$  steps
- Search the whole list for the smallest element, and then put it first
  - Find the smallest remaining thing, put it second.
  - Repeat until done.

How long does this take? Each step has to go through the whole list.

$n$  steps, go through whole list each step

$\rightarrow n \cdot n \rightarrow O(n^2)$ .

Fine for  $\approx 100k$  things, but not  $\approx 1B$

Merge Sort

$O(n \cdot \log(n))$

very slightly bigger  
than  $O(n)$