

Mon, March 27, 2023

Lecture #27

MSSC 6000

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Announcements

* Normal Office Hours today

1pm - 2pm
in CU 307

* OHI Wed are moved to 4:30pm - 5:30pm (Teams)

* HW 4 due Mon, Apr 3.

* Fri, Apr. 7 - no class

Mon, Apr. 10 - no lecture (home work day)

no OHI

Topic 11 - Hill Climbing (continued)

Inspired by Gradient Ascent:

MH #2: Steepest Ascent Hill Climbing

(for discrete only)

$x = \text{random element of } S$

(2)

while True:

$N = \text{nbhd}(x)$

$s = \text{element of } N \text{ with the highest score}$

if $\text{score}(s) > \text{score}(x)$:

$x = s$

else:

we're at the top of a hill
quit

If continuous, N is probably an infinite set, so we can't compute the score of everything in N .

Pros

* Guaranteed to find a local optimum

Cons

* Unlikely to find global opt. unless lucky and/or search space is nice

* very slow because (3)
the neighborhoods can
be big and we are
forced to score every thing
in the nbhd.

What's the slow part?

Only doing two things:

- (1) generating the nbhd
- (2) scoring the nbhd

TSP - scoring a tour of 300 cities
is not too bad

300 distance calculations

$$d((x_1, y_1), (x_2, y_2)) = \sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2}$$

two subtractions
two squarings
one addition
one square root

But the size of
the nbhd is

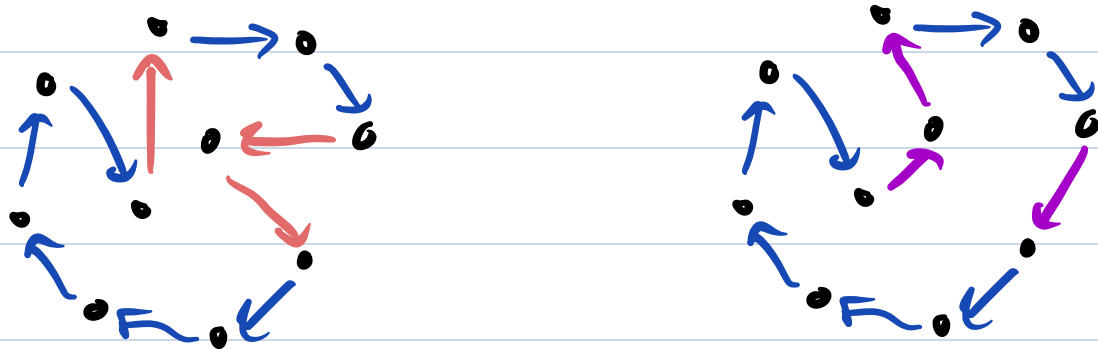
$$\binom{299}{2} = 44,551$$

Doing 300.44551 distance calculations is (4)
slow.

How can we speed this up?

(1) When you want the score of a tweaked tour, start with the score of the original tour, and alter it accordingly.

Original Tour \longrightarrow Nbrhd



Score of new = [Score of old] - [3 orange distances] + [3 purple dist]

(2) Pre-calculate and store the distance between all $\binom{300}{2} = 150 \cdot 299$ pairs of cities.

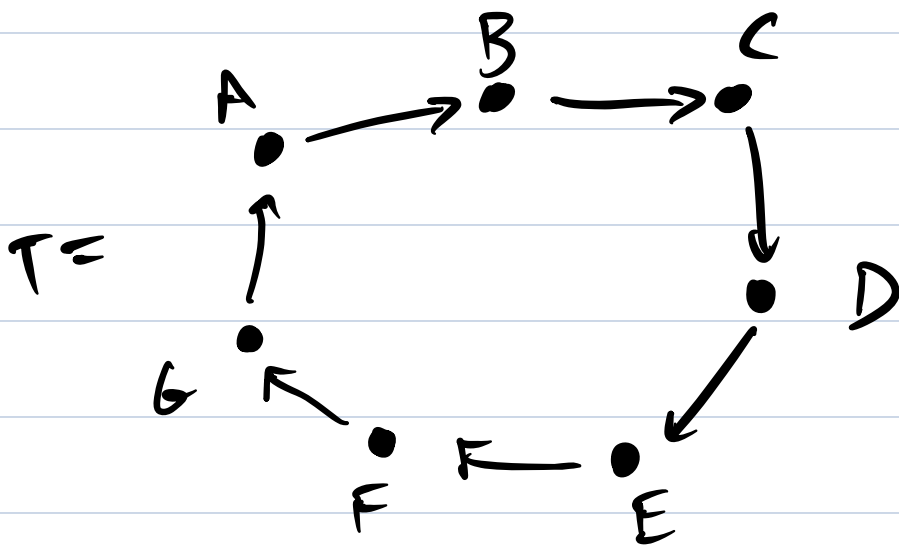
Now you don't need to do any more distance calculations. (5)

You do still need to add all these distances for each tour.

For large problems, may not have enough memory.

Can be combined with (1).

Concrete example of (1):



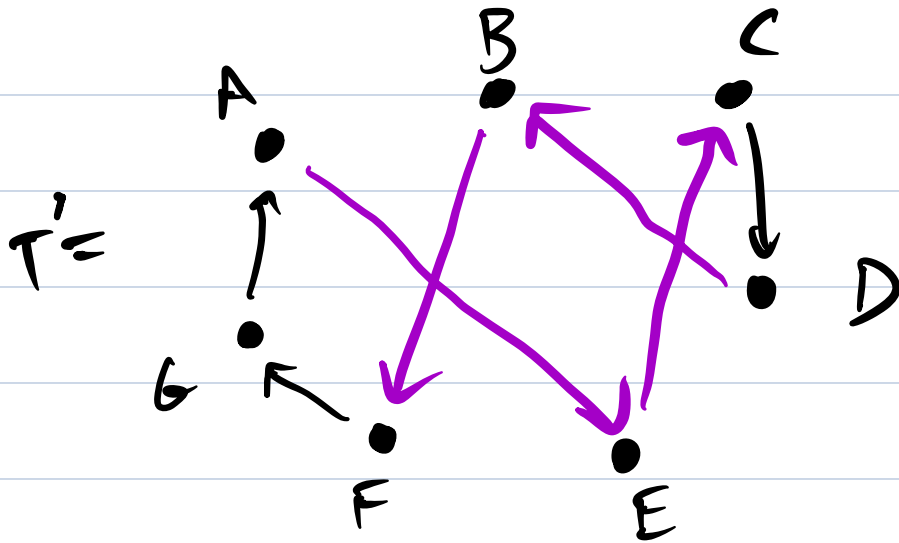
Let d be the distance function.

$$\text{score}(T) = d(A,B) + d(B,C) + d(C,D) + d(D,E) + d(E,F) + d(F,G) + d(G,A)$$

Swap B + E

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow A$ (6)

$A \rightarrow E \rightarrow C \rightarrow D \rightarrow B \rightarrow F \rightarrow G \rightarrow A$



$$\text{score}(T') = \text{score}(T) - [d(A,B) + d(B,C)$$

$$+ d(D,E) + d(E,F)] + [d(A,E)$$

$$+ d(E,C) + d(D,B) + d(B,F)]$$

If you have 300 cities:

new = old - 4 edges + 4 edges

8 distance calculations vs. 300

$$\frac{300}{8} = 37.5 \times \text{faster} \quad (7)$$

Is this tweak (swapping two cities)
a good tweak?

* makes small changes

* gets good results

Demo, 50 cities, gives a bad
result

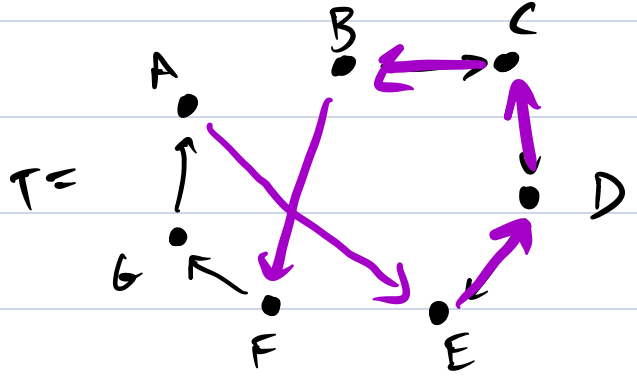
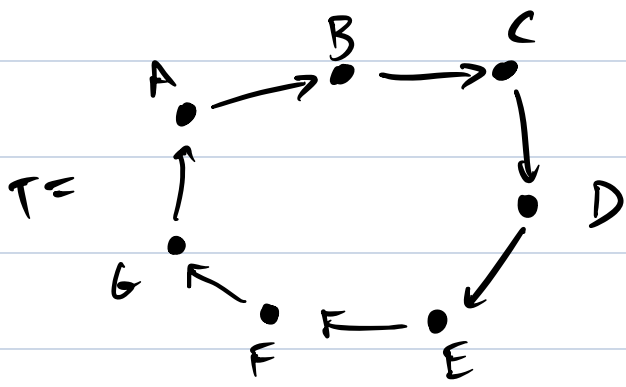
Can we think of another tweak?

Pick two cities and reverse the
entire route in between them.

A → B → C → D → E → F → G → A

A → E → D → C → B → F → G → A

8



changes 2 edges at a time, not 4

Many other possible tweaks.

Demo - much faster when we use the

scoring tricks

- this new tweak (reverse a block)

gives better results than the

old one

- still slow because it looks at the

whole nbhd

How can we adapt this for continuous spaces (when the nbhd is infinite?)

MH #3 n-trial steepest ascent

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$x = \text{random element of } S$

while True:

temp = x

repeat n times:

s = tweak(x)

if score(s) > score(temp):

temp = s

x = temp

(temp is the best of n tweaks)

if none of the tweaks beat x, then it stays the same

tweak = a random thing in the nbhd, and there are many different ways to do that

Next time: the $n=1$ version, which is just called "hill climbing".