

Fri, March 24, 2023

Lecture #26

MSSC 6000

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Announcements

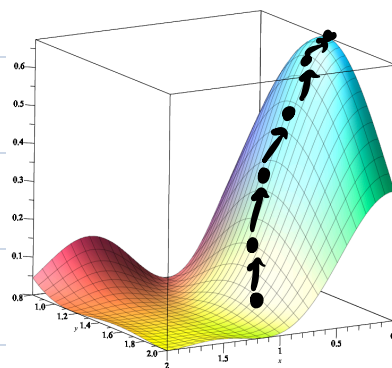
- * Normal Office Hours Monday
- * OH Wed will be moved
- * HW 4 assigned today, due Mon, Apr 3
 - ↳ on D&L, two attachments

Topic 10 - Introduction to Metaheuristics

Gradient Ascent:

- * start at a point
- * compute gradient
- * move a little in that direction
- * repeat

(vector that tells you "direction of steepest ascent")



you end up at the top of some hill (2)
and if you're lucky it's the tallest hill

Adapting to a discrete search space:

Ex: TSP

- search space: all tours on the set of cities

Each tour is a location on the landscape (a place in the mountains)

- need a definition for two tours to be "nearby"

Ex: cities = $\{1, 2, 3, 4, 5\}$

tour = $3 \rightarrow 5 \rightarrow 2 \rightarrow 1 \rightarrow 4 \rightarrow 3$

what tours are "close"?

Up to you how to define this.

One possibility: the nearby tours

one the ones you get by
swapping any two [^] cities
internal

3

3 → 5 → 2 → 1 → 4 → 3



3 → 1 → 2 → 5 → 4 → 3

- * start at a random tour
- * calculate the score of all nearby tours
- * move to the best (cheapest) one
- * repeat

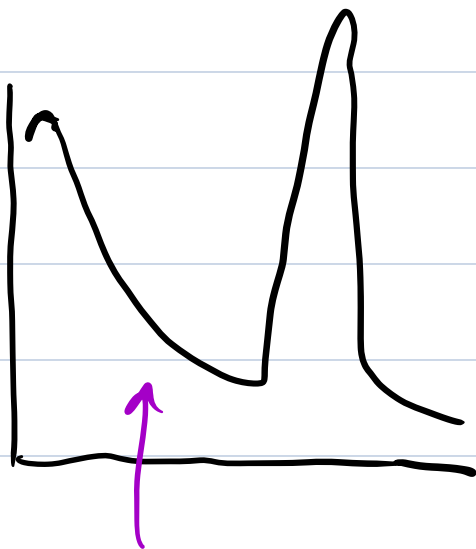
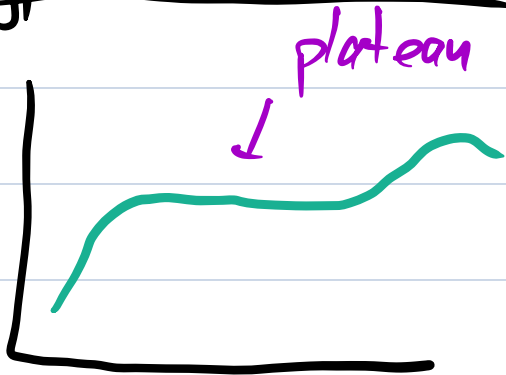
This will find a local optimum

↳ max or min
↳ best in it's area,
but maybe not best
overall

Metaheuristics are all about smart ways to explore the search space.

(4)

Typical Obstacles:



leads you the wrong way

noisy

Topic 11- Hill-Climbing

Goal: Develop a MH with Gradient Ascent as our inspiration. We want to find a global optima for either

continuous or discrete search spaces. (5)

Problem Setup:

- * Search space S full of candidates
- * Scoring function: $\text{score}(x)$, $x \in S$
(also called "fitness" or "quality")
- * A way to generate either:
 - all the candidates "nearby" a candidate (the "neighborhood of x ", $\text{nbhd}(x)$)
(doesn't make sense for continuous problems)
 - a single random candidate near a candidate
(a "tweak", $\text{tweak}(x)$)

"nearby" is up to you to define, and definitions can completely change how a MH behaves.

Two running examples:

(6)

(1) TSP

* discrete

* score = cost of tour, want to minimize

* nbhd(x) =

Suppose $x = C_1 \rightarrow C_2 \rightarrow C_3 \rightarrow \dots \rightarrow C_n \rightarrow C_1$
Define the nbhd of x to be all ways of picking two internal cities and swapping them.

How big is nbhd? $\binom{n-1}{2} = \frac{(n-1)(n-2)}{2}$
 $\approx \frac{n^2}{2}$

Pretty big nbhd.

* tweak(x): one random thing in the nbhd

(2) optimizing a continuous function in two variables $f(x, y)$

* continuous

* Search space

(7)

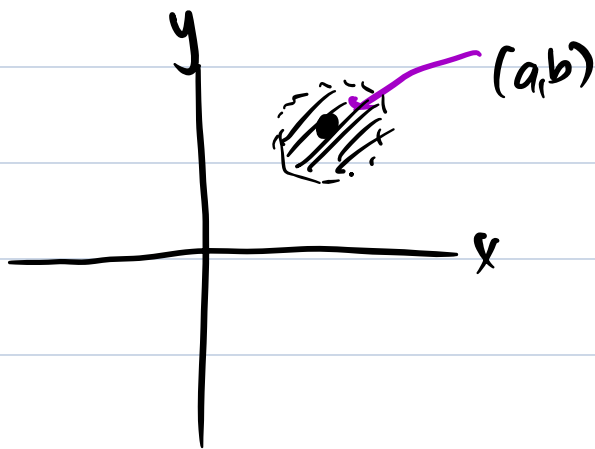
all (x,y) points maybe within
some interval

* score of a point = the value
of f at that point

* nbhd (x,y)

(what points are "near" a point
in 2D, 3D, ..., 20D, space?)

all points with a distance $< \delta$
from x for some small $\# \delta$.



* $\text{tweak}(x,y) =$
a random point in
the neighborhood, like
before.

MH #1: Random Search

(8)

best = random element of S

while True: (run until you're bored)

x = random element of S

if $\text{score}(x) > \text{score}(\text{best})$:

best = x

Possible stopping conditions:

* best score does not improve for N consecutive tries

* a preset # of attempts

* run until you get bored

This is not a good MH! It does not use any of the old information to guide future choices.

Inspired by Gradient Ascent:

MH #2: Steepest Ascent Hill Climbing
(for discrete only)

$x = \text{random element of } S$

⑨

while True:

$N = \text{nbhd}(x)$

$s = \text{element of } N \text{ with the highest score}$

if $\text{score}(s) > \text{score}(x)$:

$x = s$

else:

we're at the top of a hill

quit

If continuous, N is probably an infinite set, so we can't compute the score of everything in N .