

MSSC 6000

March 2 2022 - Day 17

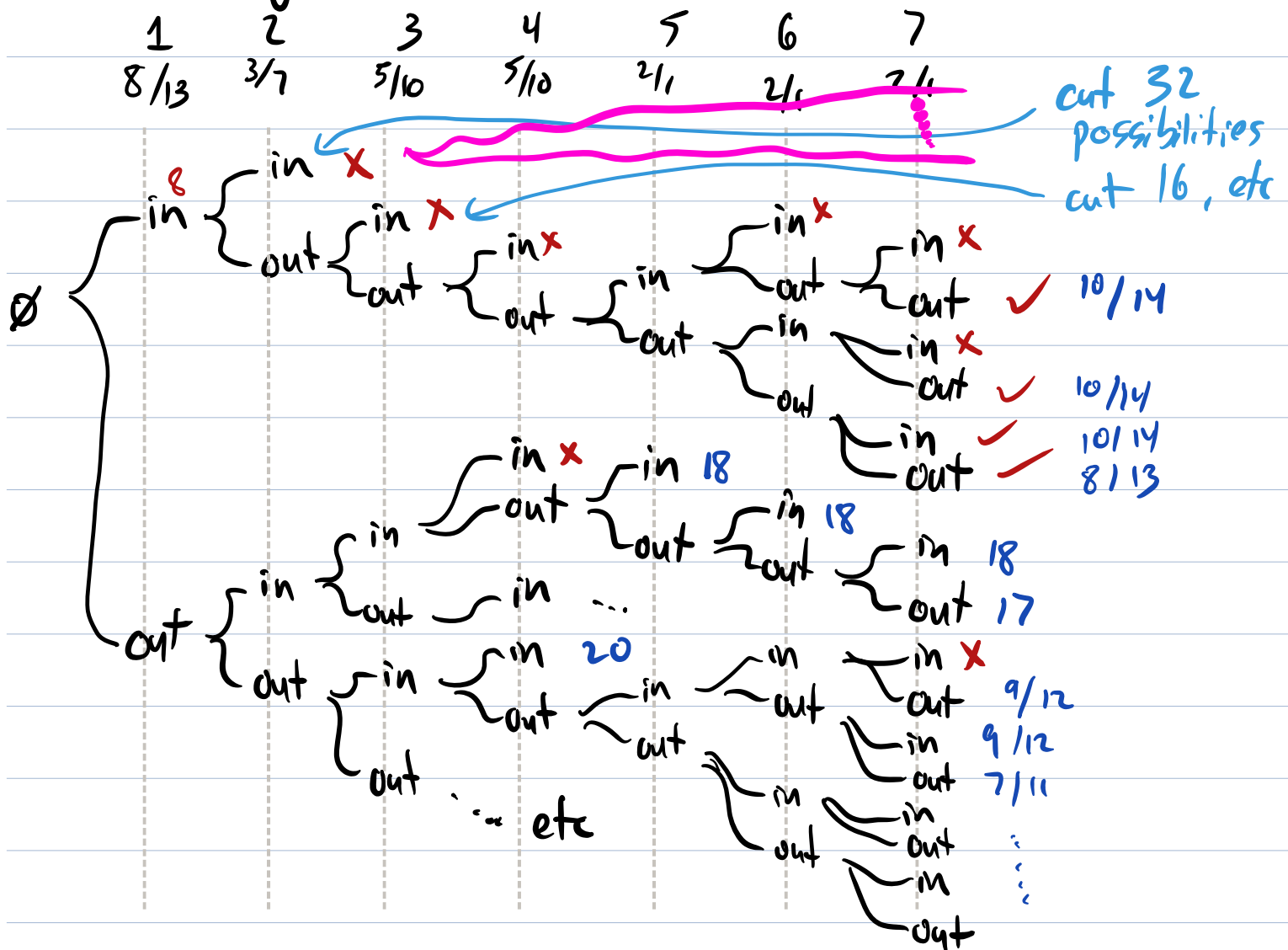
(1)

Topic #7 - Backtracking

(continued)

Backtracking:

weights / values



What are we doing?

- Putting a hierarchy on decisions that builds the whole search space,

with the critical property that if a partially built solution is bad, then every way of completing it must also be bad

Ex #2: Sudoku

1-9

$1, 9$ SAT-solver

- Start filling in blank cells left-to-right, then top-to-bottom.

- Start each blank cell with 1.

- If that cell doesn't violate a rule, move to the next cell and repeat.

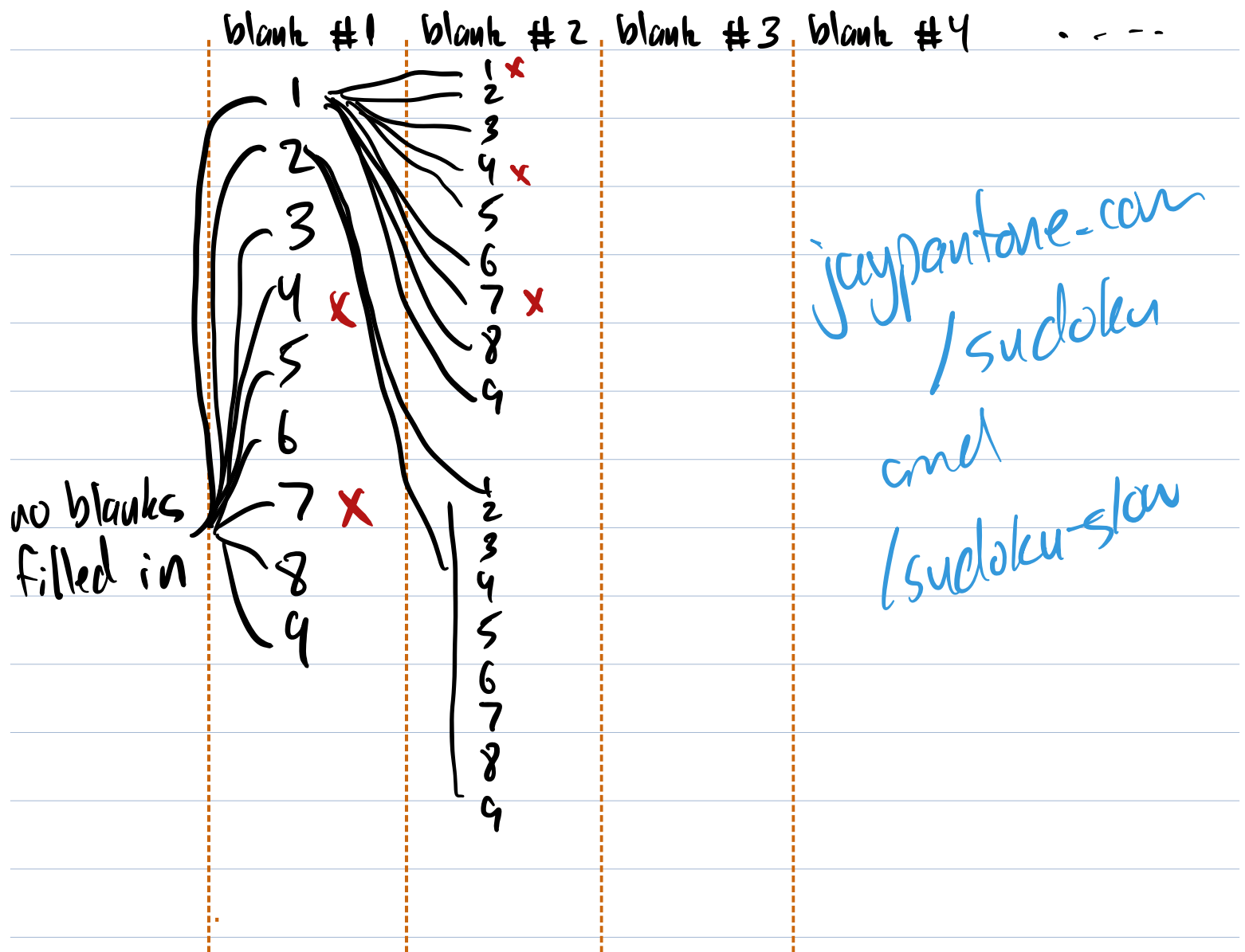
- If it does violate, increase it by 1

- If you run out of possibilities, backtrack, go back to the previous cell, increment it by 1, and repeat.

- What would a brute force algo be?
30 blanks $\Rightarrow q^{30}$ solutions to check

4	7		6	2				
6		8		5	4			
		5		(9)	8	7		4
8			4	3	2			
	3			1			4	
			9	8	7			1
1		3	8	(7)		4		
			3	4		5		9
				6	9		1	8

q^{28}



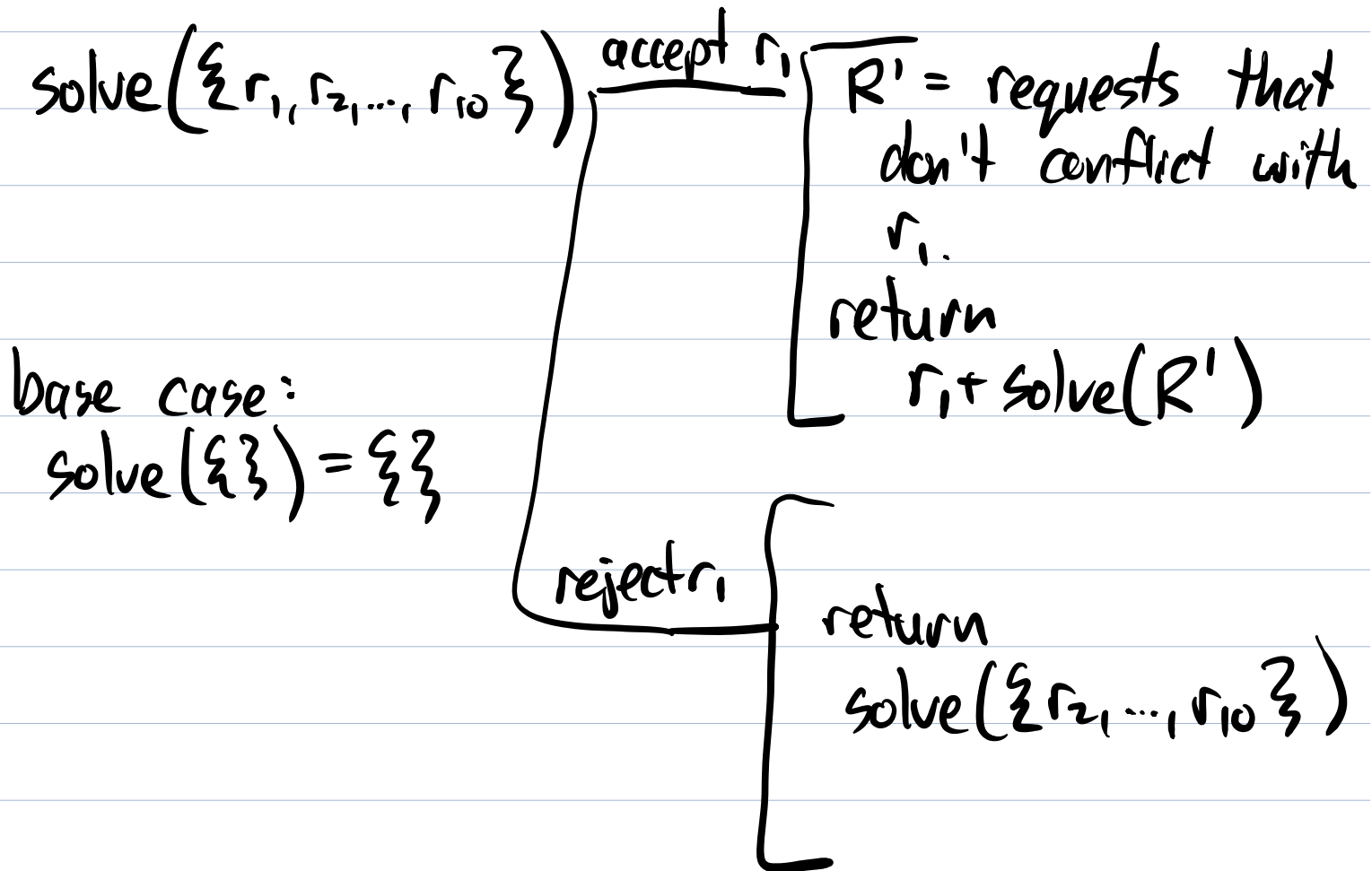
Ex 3: Weighted Interval Scheduling.

Requests $R = \{r_1, r_2, r_3, \dots\}$

- You either accept or reject each request.

If you accept r_i , then in the future you can ignore any requests that conflict with r_i .

This is what recursion is great at because we're repeating the same logic on subproblems.



Pseudo code :

function $\text{solve}(\text{requests})$:

#goal: return the best subset of requests

if $\text{len}(\text{requests}) = 0$:

return $[\]$

$\text{new_request} = \text{requests}[0]$

$\text{compatible} = \text{the requests compatible with}$

$$\text{accept_solution} = \overset{\text{new_request}}{[\text{new_request}]} + \text{solve}(\text{compatible})$$
$$\text{reject_solution} = \text{solve}(\text{requests}[1:])$$

return whichever of accept_solution and reject_solution has the higher score

Jupyter notebook demo