

MSSC 6000

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Feb 2, 2022

Announcements:

→ HW 1 due next Wed on D2L

## Lecture 3 - Greedy Algorithms

A lot of topics in this class are in the category of "problem solving paradigms".

A catalogue of ways to approach new problems. "heuristics"

What is a problem?

You have input data and/or constraints.

One question might be:

- Is it possible to satisfy all of the constraints?

Ex: Every year, the NFL has to come up with a season schedule.

There are many constraints! (2)

- 32 teams in 2 conferences
- Each conference is split into 4 divisions of 4 teams each.

Each team plays:

- all 3 division rivals, twice  
one home / one away
- each team in some other division in the same conference  
two home / two away
- four teams in the other conference  
two home / two away
- two more teams in their own conference IH / IA

- stadium constraints
- TV constraints
- holiday games
- bye week

Q: Can this be done?

A: Yes. The NFL says they use 100s of computers. They come up with ~1000 good schedules. Humans pick the best out of those.

- Another type of question:

(3)

Which solution is optimal?

(lowest cost, highest value, etc)

Ex: If Amazon has 100 packages to deliver to different houses in Milwaukee, and 5 delivery vans, which route

- uses the least gas?
- travels the fewest miles?
- takes the least time?

or some combination of these

## Greedy Algorithms

Vague Definition: A greedy algorithm is a way of solving a problem that builds up a solution bit by bit, always picks the next bit that is the best, even if that leads to a suboptimal solution in the end.

They are:

- normally lightning fast
- much better than random solutions

- sometimes pretty bad,  
sometimes pretty good,  
sometimes provably optimal,  
depending on the problem.

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Ex: Giving change - How does a cashier give change? Suppose you owe \$3.27 and pay with \$20. They start giving you bills and coins from largest to smallest.

$$\$20 - \$3.27 = \$16.73$$

<del>\$100</del>	<del>\$50</del>	<del>\$20</del>	\$10	\$5	\$1	\$0.25	\$0.10	<del>\$0.05</del>	
			1	1	1	2	2		\$0.01
			6.73	1.73	0.73	0.23	0.03		3
									0

→ 1 \$10, 1 \$5, 1 \$1, 2Qs, 2Ds, 3Ps

= 10 things. This is a greedy algorithm.

Is it optimal?

fewest # of bills/coins

Theorem: For the US currency denoms listed above, the cashier's algorithm is optimal.

Proof: To simplify things, let's assume  
denoms of 1, 5, 10, 25 cents.  
Suppose we are making  $x$  cents.

(5)

Lemma: The optimal solution will  
have  $\leq 4$  pennies.

Proof: If the sol. had  $\geq 5$  P's, we  
could replace 5P with 1 N and  
get a better solution.

Lemma: The optimal solution will  
have  $\leq 1$  N.

(Same proof)

Lemma: The optimal sol. will have  
 $\#N + \#D \leq 2$

Proof: 2N: bad

1N + 2D: bad  $\rightarrow$  1Q

0N + 3D: bad  $\rightsquigarrow$  1N + 1Q

Main proof, by induction:

Base case: 0¢ sol: 0 coins optimal ✓

Now assume we're making  $x$ ¢ worth  
of change.

Case 1:  $x < 5$ ,  $x$  pennies, optimal ✓ (6)

Case 2:  $5 \leq x < 10$ , must have 1N  
(because the lemma says no more than 4 pennies)

1N + (greedy sol. for  $x-5$  cents)  
optimal by induction ✓

Case 3:  $10 \leq x < 25$ , similar to Case 2

Case 4:  $x \geq 25$ : must have 1Q  
(or else  $> 4$  P or  $> 1$  N or  $1N+2D$ ,  
or  $3D$ )

sol: 1Q + (greedy sol. for  $x-25$  cents). 

Is the cashier's algo. optimal for any set of denominations? No.

Ex: US Postage Denominations

1, 2, 3, 5, 10, 20, 35, 36, 55, 65, 75, 95, 100,  
120, 200, 500, 795, 1000, 2635.

To make 72¢: 65 + 5 + 2      3 stamps

Better: 36 + 36 = 72      2 stamps

Greedy  $\neq$  Optimal

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Throughout the course we're going to learn about a catalogue of problems that model all kinds of real world problems that you might face.

Problem #1: Interval Scheduling